# Population Games for Cognitive Radios: Evolution through Imitation

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### **Outlines**

- Introduction: evolutionary game theory and population games
- Population games applied to CR: model and assumptions
  - Simple and double imitation w/o channel constraint
  - Simple and double imitation with channel constraint
- Conclusion and future work

### **Evolutionary games overview**

- Evolutionary games formalism is a central mathematical tool developed by biologists for predicting populations dynamics in the context of interactions between populations.
- All players in a population are programmed to use strategies
- Strategies with high payoff will spread within the population. This can be achieved by learning, copying or inheriting strategies.
- The payoff depends on the frequency of the strategies within the population. Since this frequencies change according to the payoffs, this yields a feedback loop.

- Population games are a particular class of evolutionary games
- They model strategic interactions in which [1]
  - Population is large
  - The number of strategies is finite
  - Agents interact at random (e.g. pairwise)
  - Payoffs are continuous

#### **Definitions**

- Strategies:  $S = \{1, ..., n\}$
- Population states:  $X = \{x \in \mathbb{R}^n_+, \sum_{i=1}^n x_i = 1\}$
- Payoff function:  $\pi: X \to R^n$  assigns to each state a payoff vector
- Payoff component for strategy  $i: \pi_i: X \to R$
- Average payoff in state x:  $\bar{\pi}(x) = \sum_i x_i \pi_i(x)$

### Evolutionary game dynamics

- Players play mixed strategies
- Players update their strategy according to their environment
- The updating process is called a **revision protocol**: Let  $\rho_{i,j}: R^n \times X \to R_+$  be the switch rate from i to j i.e.,  $\rho_{i,j}dt$  is the probability for a player to switch from i to j in dt.
- The revision protocol generates a system dynamic (Kolmogorov):

$$x_i(t+dt) = x_i(t) + \sum_{j \in \mathcal{S}} x_j(t) \rho_{j,i} dt - x_i(t) \sum_{j \in \mathcal{S}} \rho_{i,j} dt$$

$$\dot{x}_i(t) = \sum_{j \in \mathcal{S}} x_j(t) \rho_{j,i} - x_i(t) \sum_{j \in \mathcal{S}} \rho_{i,j}$$



Example of revision protocol: the proportional imitation rule (PIR)

- A revising player adopting strategy i picks at random an opponent
- ullet It observes its current strategy j and payoff  $\pi_j$
- And switches to j iff  $\pi_j > \pi_i$  with a probability proportional to the payoff difference
- The switch rate is:  $\rho_{i,j} = x_j \sigma [\pi_j \pi_i]_+ (\sigma \text{ is a constant})$
- The resulting system dynamic is the replicator dynamic [2]:

$$\dot{x}_i(t) = \sigma x_i(t)(\pi_i(t) - \bar{\pi}(t))$$



Another revision protocol: the double imitation (DI)

- A revising player adopting strategy *i* picks at random two opponents
- ullet It observes their current strategy  $j_1$  and  $j_2$  and payoffs  $\pi_{j_1}$  and  $\pi_{j_2}$
- The switch rate is a function of  $\pi_i$ ,  $\pi_{j_1}$ ,  $\pi_{j_2}$  and two control parameters  $(\alpha, \omega)$
- The system dynamic is the aggregate monotone dynamic [2]:

$$\dot{x}_i(t) = rac{x_i(t)}{\omega - lpha} \left[ 1 + rac{\omega - ar{\pi}(t)}{\omega - lpha} 
ight] (\pi_i(t) - ar{\pi}(t))$$



# Population Games applied to CR

Game G and cognitive radio (CR) scenario

- There are N SUs and C channels, each with availability  $\mu_i \in [0,1]$
- Strategies:  $S = \{1, ..., C\}$
- Throughput on channel i (normalized): RV  $T_i$
- Payoffs: expectation of the normalized throughput

### Model assumptions:

- Imitation can be performed across channels
- Generic MAC protocol. The throughput of the SUs on the same channel i is defined as:

$$\pi_i(\mu_i, x_i) = E[T_i] \approx \mu_i/n_i = \mu_i/(x_iN)$$

Payoffs are obtained at the end of each iteration without errors

#### **Theorem**

In the asymptotic case where N is large, G admits a unique NE. At the NE, there are  $x_i^*$  N SUs staying on channel i, where  $x_i^* = \frac{\mu_i}{\sum_{l \in C} \mu_l}$ .

• G is a congestion game and also a potential game

$$P(\mathbf{x}) \triangleq \sum_{i \in \mathcal{C}} \int_{\epsilon_0}^{x_i N} \frac{\mu_i}{t} dt$$
 and  $\frac{\partial P(x)}{\partial x_i} = \mathbb{E}[\pi_j(\mu_i, x_i)]$ 

- The problem  $\max_{\mathbf{x}} P(\mathbf{x}) \ s.t. \sum_{i \in \mathcal{C}} x_i = 1$  has a unique solution
- Convergence is exponential (DI converges at a higher rate)



Distributed algorithm based on imitation Assumptions:

- ullet Current payoff  $\pi_i$  is included in the header of each transmitted packet
- Each SU is able to overhear one or two packets of other SUs

 $\begin{array}{lll} \textbf{Algorithm 1} & \textbf{Imitation Spectrum Access Policy (ISAP) executed at each} \\ \textbf{Secondary User} & \end{array}$ 

- 1: Initialization: Set  $\epsilon_t$
- 2: At each iteration t
- 3: With probability  $1 \epsilon_t$  perform imitation (PIR or DI)
- 4: With probability  $\epsilon_t$  switch to a random channel

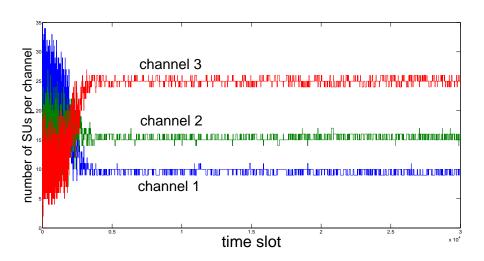


Figure: PIR-ISAP: number of SUs per channel as a function of time

- SUs can overhear only on the channel on which they stay
- They imitate the payoff obtained at time t-1.

### Theorem (Dynamics)

In the case of proportional imitation policy it holds that:

$$x_i(t+1) = \sum_{j,l,k \in \mathcal{C}} \frac{x_j^l(t)x_j^k(t)}{x_j(t)} F_{l,k}^i \quad \forall i \in \mathcal{C}$$

Differently, the double imitation policy yields:

$$x_i(t+1) = \sum_{j,l,k,z \in \mathcal{C}} rac{x_j^l(t)x_j^k(t)x_j^z(t)}{[x_j(t)]^2} F_{l,\{k,z\}}^i \quad \forall i \in \mathcal{C}$$

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The system dynamics are well approximated by a double replicator dynamic (PISAP) and by a double aggregate monotone dynamic (DISAP):

• Double replicator, e.g., has the following expression:

$$\begin{cases} x_i(u) = x_i(u-1) + \sigma x_i(u-1)[\pi_i(u-1) - \bar{\pi}(u-1)] \\ x_i(v) = x_i(v-1) + \sigma x_i(v-1)[\pi_i(v-1) - \bar{\pi}(v-1)] \end{cases}$$

where u = 2t, v = 2t + 1.

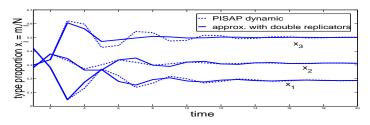


Figure: System dynamic and its approximation by double replicator dynamic.

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### Finite populations:

- p denotes the population proportions whenever the population is finite.
- It holds that  $p \rightarrow x$  when N is very large.

### Theorem (Finite populations)

For any imitation rule F, if the imitation among SUs of the same type occurs randomly and independently, then  $\forall \delta > 0$ ,  $\epsilon > 0$  and any initial state  $\{\widetilde{x}_i(0)\}$ ,  $\{\widetilde{x}_i(1)\}$ , there exists  $N_0 \in \mathbb{N}$  such that if  $N > N_0$ ,  $\forall i \in \mathcal{C}$ , the event  $|p_i(t) - x_i(t)| > \delta$  occurs with probability less than  $\epsilon$ , where  $p_i(0) = x_i(0) = \widetilde{x}_i(0)$ ,  $p_i(1) = x_i(1) = \widetilde{x}_i(1)$ .

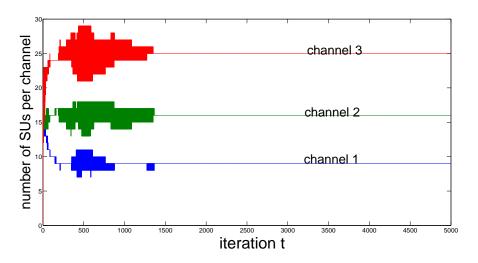


Figure: DISAP: number of SUs per channel as a function of time with channel constraint.

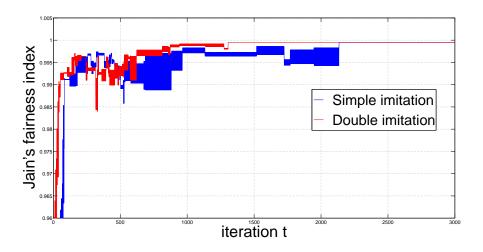


Figure: PISAP and DISAP fairness trends comparison (one realization).

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### Conclusion and future work

- Imitation-based Spectrum Access Policies allow the SUs to load-balance the system throughput.
- The approach is totally distributed and relies solely on local interactions amongst users
- Our next goals are to make the model more realistic and adapt our algorithms accordingly
  - More realistic MAC protocol (CSMA/CA)
  - Different imitation strategies
  - Non-symmetric topologies
  - Priority schemas

### References I

- [1] W. H. Sandholm. Local Stability under Evolutionary Game Dynamics. Theoretical Economics, 5, 2010.
- [2] K. H. Schlag. Why Imitate, and if so, How? Discussion paper, University of Bonn, Feb. 1996.