
Incremental Spectral Clustering with the Normalised Laplacian

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Abstract

Partitioning a graph into groups of vertices such that those within each group are more densely connected than vertices assigned to different groups, known as *graph clustering*, is often used to gain insight into the organization of large scale networks and for visualization purposes. Whereas a large number of dedicated techniques have been recently proposed for static graphs, the design of *on-line* graph clustering methods tailored for evolving networks is a challenging problem, and much less documented in the literature. Motivated by the broad variety of applications concerned, ranging from the study of biological networks to graphs of scientific references through to the exploration of communications networks such as the world wide web, it is the main purpose of this paper to introduce a novel, computationally efficient, approach to graph clustering in the evolutionary context. Namely, the method promoted in this article is an incremental eigenvalue solution for the spectral clustering method described by Ng. *et al.* (2001). Beyond a precise description of its practical implementation and an evaluation of its complexity, its performance is illustrated through numerical experiments, based on datasets modelling the evolution of a HIV epidemic and the purchase history graph of an e-commerce website.

1 Introduction

Graph-mining has recently received increasing attention in the machine-learning literature, motivated by application domains such as the Internet, social networks, epidemics of transmissible infectious diseases, sensor and biological networks. A crucial task in exploratory analysis and data visualization is *graph clustering* [1, 2], which aims to partition the vertices in a graph into groups or clusters, with dense internal connections and few connections between each other. There is a large body of work on graph clustering. A possible approach is to consider a certain measure that quantifies community structure and formulate the clustering issue as an optimization problem (which is generally NP-hard), for which fairly good solutions can be obtained recursively or by using adequate metaheuristics, see [3, 4, 5, 6] for example. In this paper, focus is on the *spectral clustering* approach [7], which performs well empirically, is often simple to implement and benefits computationally from fast linear algebra libraries. The general principle of spectral clustering is to compute the smallest eigenvectors of some particular matrix \mathbf{L} (refer to Section 2 for further details) and then cluster the vertices based on their representation in the corresponding eigen-space. The popularity of this approach arises from the fact that the obtained clustering of vertices is closely connected to the spectral relaxation of the minimization of the normalized cut criterion, see [8].

In many applications such as communications networks (e.g. the Web and Internet), biological networks (of proteins, metabolic reactions, *etc.*), social networks or networks of scientific citations for instance, the graphs of interest slowly change with time. A naive approach to this incremental problem is to cluster each graph in the sequence separately, however this is computationally expensive for spectral clustering as the cost of solving the eigenvalue problem is $O(n^3)$ at each iteration,

where n is the number of vertices. There has been some previous work on the incremental spectral clustering problem, for example [9, 10, 11] however only [10, 11] update the eigen-system. In this paper we propose an efficient method for clustering a sequence of graphs which leverages the eigen-decomposition and the clustering on the previous graph to find the new clustering. Firstly, a fast approximation of a rank- k eigen-decomposition of \mathbf{L}_{t+1} knowing that of \mathbf{L}_t is derived from the Singular Value Decomposition (SVD) updating approach used for Latent Semantic Indexing in [12]. Here, the update is efficient to compute when the change (that will be described more precisely in the subsequent analysis) between \mathbf{L}_t and \mathbf{L}_{t+1} is a low-rank matrix. Secondly the clustering of vertices is updated accordingly to the new eigen-space. The efficiency of the complete approach, in term of clustering accuracy, is demonstrated on real data.

The paper is organised as follows. A Laplacian-based graph clustering approach is recalled in Section 2. Then Section 3 details the proposed eigen-decomposition update and Section 4 describes its use in the framework of spectral clustering. Numerical results are gathered in Section 5, and the paper ends with Section 6 discussing results and planned future work.

Notation: A bold uppercase letter represents a matrix, e.g. \mathbf{X} , and a column vector is displayed using a bold lowercase letter, e. g. \mathbf{x} . The transpose of a matrix or vector is written \mathbf{X}^T . The concatenation of the columns of two matrices \mathbf{A} and \mathbf{B} is written $[\mathbf{A} \ \mathbf{B}]$.

2 Laplacian Based Graph Clustering

Consider an undirected graph $G = (V, E)$, composed of a set of vertices $\{v_1, \dots, v_n\} = V$ and edges $E \subseteq V \times V$ such that for every edge (v_i, v_j) there is also an edge (v_j, v_i) . One way of representing the edges of G is using an *adjacency matrix* \mathbf{A} which has $\mathbf{A}_{ij} = 1$ if there is an edge from the i -th to the j -th vertex otherwise $\mathbf{A}_{ij} = 0$. More generally, the *weight matrix* \mathbf{W} allows one to assign nonzero numerical values on edges, and thus $\mathbf{W}_{ij} \neq 0$ when there is an edge from the i -th to the j -th vertex.

In the perspective of spectral clustering, a useful way of representing G is through its Laplacian matrix [13]. There are several definitions of the Laplacian matrix, however we are interested in the *normalised Laplacian matrix*, which is symmetric and positive semi-definite:

Definition 2.1. *The unnormalised Laplacian matrix of a graph G is defined as $\mathbf{L} = \mathbf{D} - \mathbf{W}$ where \mathbf{D} is the degree matrix with diagonal entries $\mathbf{D}_{ii} = \deg(v_i)$, denoting by $\deg(v_i) = \sum_j \mathbf{W}_{ij}$ the degree of the i -th vertex, and zeros elsewhere, and \mathbf{W} is the weight matrix. The normalised Laplacian matrix of a graph G is then defined as $\tilde{\mathbf{L}} = \mathbf{D}^{-\frac{1}{2}} \mathbf{L} \mathbf{D}^{-\frac{1}{2}}$.*

The normalised Laplacian matrix is used in the spectral clustering approach of Ng et al. [14]. The algorithm computes the Laplacian and then finds the k smallest eigenvectors which form the columns of a matrix \mathbf{V} . This matrix is then normalised so that its rows have unit norm, which are then clustered with the k -means algorithm. In general the Laplacian matrix is often sparse and one can use power or Krylov subspace methods such as the Lanczos method to find the eigenvectors. There are several variants of this algorithm, such as that of [8] which uses the so-called random walk Laplacian and clusters the smallest eigenvectors in a similar way. One of the motivations for these clustering methods is from the spectral relaxation of the minimisation of the normalised cut criterion.

A naive approach to spectral clustering on a sequence of graphs has a large update cost due to the computation of the eigen-decomposition of $\tilde{\mathbf{L}}_t$ at each iteration t . A more efficient approach is to update relevant eigenvectors from one iteration to the next. In [9] the authors use the spectral clustering of [14] however they do not update the eigenvectors incrementally but instead the clustering directly. The iterative clustering approach of Ning et al. [10, 11] uses the spectral clustering method given in [8] and updates eigenvectors incrementally. The algorithm incrementally updates the generalised eigenproblem $\mathbf{L}\mathbf{v} = \lambda\mathbf{D}\mathbf{v}$ by finding the derivatives on the eigenvalues/vectors with perturbations in all of the quantities involved. An iterative refinement algorithm is given for the eigenvalues and eigenvectors given a change in the edges or vertices of a graph. One then clusters the resulting k smallest eigenvectors using k -means clustering. In order to limit errors which can build up cumulatively, the authors recompute the eigenvectors after every R -th graph in the sequence.

A disadvantage of the approach of [10, 11] is that to update an eigenvector one must invert a small matrix for each weight change which makes updates costly. The size of this matrix is proportional to the number of neighbours of the vertices incident to the changed edge.

3 Incremental Eigen-approximation

In this section we address the updating problem upon the rank- k eigen-decomposition of a symmetric matrix, *i.e.* we are interested only in the largest k eigenvalues/vectors of the corresponding matrix: $(\mathbf{A}^T \mathbf{A})_k + \mathbf{U}$, where $\mathbf{U} = \mathbf{Y}_1 \mathbf{Y}_2^T + \mathbf{Y}_2 \mathbf{Y}_1^T$ and $(\mathbf{A}^T \mathbf{A})_k$ is the rank- k eigen-decomposition of $\mathbf{A}^T \mathbf{A}$.¹ This update is a general operation, however it will be explained in the context of spectral clustering later in Section 4. We assume that $\mathbf{Y}_1, \mathbf{Y}_2 \in \mathbb{R}^{n \times p}$ and one does not have direct access to $\mathbf{A} \in \mathbb{R}^{m \times n}$ but only the matrix $\mathbf{A}^T \mathbf{A}$.

To compute this update, we first project the columns of \mathbf{Y}_1 into the space orthogonal to the eigenvectors \mathbf{Q}_k of $(\mathbf{A}^T \mathbf{A})_k$, a process known as *deflation*. Hence the matrix \mathbf{Y}_1 is deflated as follows: $\bar{\mathbf{Y}}_1 = (\mathbf{I} - \mathbf{Q}_k \mathbf{Q}_k^T) \mathbf{Y}_1$. Note that $\bar{\mathbf{Y}}_1 \Theta_1$ for some Θ_1 , is orthogonal to \mathbf{Q}_k since $\mathbf{Q}_k^T \bar{\mathbf{Y}}_1 \Theta_1 = (\mathbf{Q}_k^T \mathbf{Y}_1 - \mathbf{Q}_k^T \mathbf{Y}_1) \Theta_1 = \mathbf{0}$. If we take the SVD $\bar{\mathbf{Y}}_1 = \bar{\mathbf{P}}_1 \bar{\Sigma}_1 \bar{\mathbf{Q}}_1^T$ then $\bar{\mathbf{P}}_1$ is orthogonal to \mathbf{Q}_k since $\bar{\mathbf{P}}_1 = \bar{\mathbf{Y}}_1 \bar{\mathbf{Q}}_1 \bar{\Sigma}_1^{-T}$ assuming $\bar{\Sigma}_1$ has nonzero diagonal entries.

At the next stage we would like to orthogonalise the columns of \mathbf{Y}_2 with respect to both \mathbf{Q}_k and $\bar{\mathbf{P}}_1$. Hence we deflate \mathbf{Y}_2 in the following way: $\bar{\mathbf{Y}}_2 = (\mathbf{I} - \bar{\mathbf{P}}_1 \bar{\mathbf{P}}_1^T - \mathbf{Q}_k \mathbf{Q}_k^T) \mathbf{Y}_2$ where we have used the fact that $\bar{\mathbf{P}}_1$ is orthogonal to \mathbf{Q}_k . Proved in a similar way to the step used earlier, the matrix in the column space of $\bar{\mathbf{Y}}_2$, $\bar{\mathbf{Y}}_2 \Theta_2$ for some Θ_2 , is orthogonal to \mathbf{Q}_k and $\bar{\mathbf{P}}_1$. Hence, we compute the SVD $\bar{\mathbf{Y}}_2 = \bar{\mathbf{P}}_2 \bar{\Sigma}_2 \bar{\mathbf{Q}}_2^T$ and note that the matrices $\bar{\mathbf{P}}_1, \bar{\mathbf{P}}_2$ and \mathbf{Q}_k are mutually orthogonal and span the space spanned by $(\mathbf{A}^T \mathbf{A})_k + \mathbf{U}$. This allows one to write $(\mathbf{A}^T \mathbf{A})_k + \mathbf{U} = \tilde{\mathbf{Q}} \Delta \tilde{\mathbf{Q}}^T$ with $\tilde{\mathbf{Q}} = [\mathbf{Q}_k \bar{\mathbf{P}}_1 \bar{\mathbf{P}}_2] \in \mathbb{R}^{n \times (k+2p)}$ and

$$\Delta = \begin{bmatrix} \Omega_k + \mathbf{Q}_k^T \mathbf{U} \mathbf{Q}_k & \mathbf{Q}_k^T \mathbf{U} \bar{\mathbf{P}}_1 & \mathbf{Q}_k^T \mathbf{Y}_1 \bar{\mathbf{Q}}_2 \bar{\Sigma}_2 \\ \bar{\mathbf{P}}_1^T \mathbf{U} \mathbf{Q}_k & \bar{\mathbf{P}}_1^T \mathbf{U} \bar{\mathbf{P}}_1 & \bar{\Sigma}_1 \bar{\mathbf{Q}}_1^T \bar{\mathbf{Q}}_2 \bar{\Sigma}_2 \\ \bar{\Sigma}_2 \bar{\mathbf{Q}}_2^T \mathbf{Y}_1^T \mathbf{Q}_k & \bar{\Sigma}_2 \bar{\mathbf{Q}}_2^T \bar{\mathbf{Q}}_1 \bar{\Sigma}_1 & \mathbf{0} \end{bmatrix},$$

where, the above expression of Δ is simply an expansion of the matrix $\tilde{\mathbf{Q}}^T ((\mathbf{A}^T \mathbf{A})_k + \mathbf{U}) \tilde{\mathbf{Q}}$ since $\tilde{\mathbf{Q}}$ is a basis of the space occupied by $((\mathbf{A}^T \mathbf{A})_k + \mathbf{U})$. We take the rank- k eigen-decomposition $\Delta_k = \mathbf{H}_k \Pi_k \mathbf{H}_k^T$ and then the final eigen-approximation is given by $(\tilde{\mathbf{Q}} \mathbf{H}_k) \Pi_k (\tilde{\mathbf{Q}} \mathbf{H}_k)^T$.

The deflation and SVD of $\bar{\mathbf{Y}}_1$ and $\bar{\mathbf{Y}}_2$ cost $\mathcal{O}(npk)$ and $\mathcal{O}(np^2)$ respectively and the eigen-decomposition of Δ is $\mathcal{O}((k+2p)^3)$ as $\Delta \in \mathbb{R}^{(k+2p) \times (k+2p)}$. In order to compute Δ one can reuse the computations of $\mathbf{Q}_k^T \mathbf{Y}_1, \mathbf{Q}_k^T \mathbf{Y}_2, \bar{\mathbf{P}}_1^T \mathbf{Y}_2$ which are used for deflations and also the matrices used for the SVD decompositions. Thus Δ is found in $\mathcal{O}(p^3 + p^2k + pk^2)$, and the overall complexity of this algorithm is $\mathcal{O}((k^2 + p^2)(p + k) + np(p + k))$. Of note here is that the dimension n of $\mathbf{A}^T \mathbf{A}$ scales the complexity in a linear fashion, however costs are cubically related to k and p . Under the specific circumstances of working with sparse matrices and updating graph clustering this cost is later shown to reduce.

Eigenupdates corresponding to the addition and deletion of rows/columns are a particular case of the considered update, in which some rows/columns have zero values before or after the update. However, these are not detailed due to space limitations.

4 Incremental Cluster Membership

We return to the eigen-problem $\tilde{\mathbf{L}} \mathbf{v} = \lambda \mathbf{v}$ in which we are interested in the eigenvectors with the smallest eigenvalues. We define the *shifted Laplacian* as $\hat{\mathbf{L}} = 2\mathbf{I} - \tilde{\mathbf{L}} = \mathbf{I} + \mathbf{D}^{-\frac{1}{2}} \mathbf{W} \mathbf{D}^{-\frac{1}{2}}$, which

¹ We have written the update in terms of a symmetric matrix $\mathbf{A}^T \mathbf{A}$ to improve notation. Note that any positive semi-definite symmetric matrix \mathbf{M} can be decomposed into the form $\mathbf{M} = \mathbf{A}^T \mathbf{A}$, where \mathbf{A} is real, for example by using an eigen-decomposition.

is positive semi-definite since $\tilde{\mathbf{L}}$ is positive semi-definite with largest eigenvalue 2. Since we are interested in the smallest eigenvectors of $\tilde{\mathbf{L}}$ they correspond exactly to the maximum eigenvectors of $\hat{\mathbf{L}}$ and we can use the eigen-update methods described above.

Putting the ingredients together allows us to outline an efficient incremental method for performing graph clustering called Incremental Approximate Spectral Clustering (IASC), see Algorithm 1. At a high level the algorithm is quite simple: at each iteration t one mostly recovers the steps of standard spectral clustering. In this case, we use the shifted Laplacian matrix $\hat{\mathbf{L}}_t$ instead of the normalised Laplacian matrix, which does not affect the clustering as both matrices have same eigen-vectors. Furthermore we use the eigen-approximation method above to approximate the largest eigenvectors of $\hat{\mathbf{L}}_t$ using the eigenvectors computed at the previous iteration, instead of (costly) recomputing the exact eigen-decomposition at each iteration. Similarly, we initialise k -means with centers at previous iteration. Finally, we store the first k_2 eigenvectors of $\hat{\mathbf{L}}_t$, with $k_2 \geq k_1$, to reduce cumulative errors introduced in the loop at the expense of increased computation.

Algorithm 1 Incremental Approximate Spectral Clustering using the Normalised Laplacian

Require: Graphs G_1, \dots, G_T of sizes n_1, \dots, n_T , no. of clusters k_1 , matrix approximation rank $k_2 \geq k_1$, eigen-decomposition recomputation step R

- 1: **for** $t = 1 \rightarrow T$ **do**
- 2: Compute the shifted Laplacian for G_t , $\hat{\mathbf{L}}_t$
- 3: **if** $i \% R == 1$ **then**
- 4: Compute rank- k_2 eigen-decomposition of $\hat{\mathbf{L}}_t$, $\mathbf{Q}_{k_2}^{(t)} \mathbf{\Omega}_{k_2}^{(t)} \mathbf{Q}_{k_2}^{(t)T}$
- 5: **else**
- 6: Use rank- k_2 eigen-decomposition of $\hat{\mathbf{L}}_{t-1}$ and approximated eigen-update of Section 3 to compute rank- k_2 eigen-decomposition of $\hat{\mathbf{L}}_t$, $\mathbf{Q}_{k_2}^{(t)} \mathbf{\Omega}_{k_2}^{(t)} \mathbf{Q}_{k_2}^{(t)T}$
- 7: **end if**
- 8: Let $\mathbf{Q}_{k_1}^{(t)}$ be the matrix of the first k_1 columns of $\mathbf{Q}_{k_2}^{(t)}$
- 9: Normalise the rows of $\mathbf{Q}_{k_1}^{(t)}$, i.e. $\mathbf{Q}_{k_1}^{(t)} \leftarrow \text{diag}(\mathbf{Q}_{k_1}^{(t)} (\mathbf{Q}_{k_1}^{(t)})^T)^{-\frac{1}{2}} \mathbf{Q}_{k_1}^{(t)}$
- 10: Use k -means on rows of $\mathbf{Q}_{k_1}^{(t)}$, using \mathbf{c}_{t-1} to find initial centroids (if not at first iteration), and store indicators $\mathbf{c}_t \in \{1, \dots, k_1\}^{n_t}$
- 11: **end for**
- 12: **return** Cluster membership vectors $\mathbf{c}_1 \in \{1, \dots, k\}^{n_1}, \dots, \mathbf{c}_T \in \{1, \dots, k\}^{n_T}$

The complexity of Algorithm 1 is dictated by the sparsity of the graphs and the extent of the change between successive graphs in the input sequence. At iteration t and step 2 one can compute $\hat{\mathbf{L}}_t$ from the weight and degree matrix at a cost of $O(n_t + |E_t|)$. In the following step, if there is a change between edges incident to vertices $S = \{v_{I_1}, \dots, v_{I_\ell}\}$ then the rows and columns corresponding to the union of the neighbours of S , $n(S)$, will change in the corresponding shifted Laplacian. In this case $p = |n(S)|$ in $\mathbf{Y}_1, \mathbf{Y}_2 \in \mathbb{R}^{n \times p}$ and if $n(S)$ is small, this update can be efficiently computed.

5 Computational Results

In this section we compare the clustering quality of: standard spectral clustering (denoted `Exact`) Ning et al.'s incremental strategy, and IASC. We use the clustering methods on two real datasets. The first one represents individuals in Cuba who have been detected as having HIV between 1986 and 2004, see [15] for precise details of the dataset. Edges in the graph represent a sexual contact between two individuals as determined using contact tracing, which is a technique used in epidemiology to limit the spread of an infection whereby contacts of an infected person are identified and diagnosed. The full sexual contact graph at the end of 2004 consists of 5389 people however we consider the largest connected component of size 2387. We examine the epidemic graph at 3 month steps and cluster using the 3 approaches starting at the point in which the graph has 500 vertices. The experiment is run using $k_1 = 25$, $k_2 = 100$ and $R = 10$ for IASC. For the method of Ning et al. we also use 25 clusters and recompute exact eigenvectors after every 10 iterations. Note that in our implementation of k -means clustering, k is an upper bound on the number of clusters.

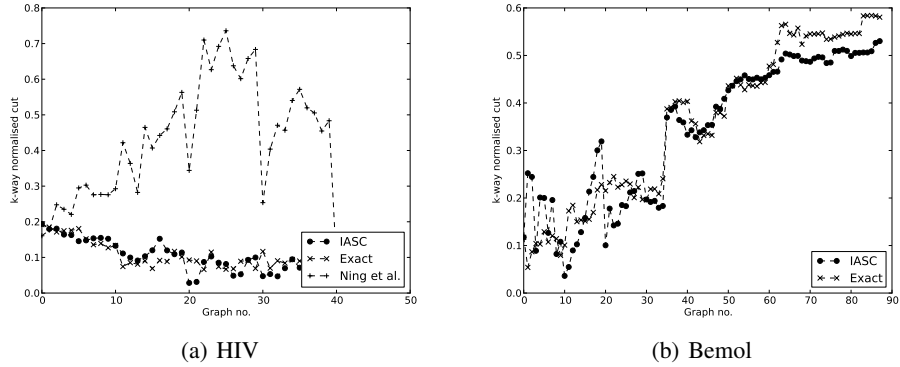


Figure 1: The k -way normalised cuts using the HIV and e-commerce graphs.

The next dataset was introduced in [16] and corresponds to the purchase history of an e-commerce website over a period of almost two years. The dataset is a bipartite graph between users and product with more than 700,000 users and 1,200,000 products. In the current experiment we focus on the first 5000 users. A graph is constructed between users with edge weights equal to the number of commonly purchased products between two users.

To evaluate the learned clusters we measure the k -way *normalised cut*, defined as follows:

$$N = \frac{1}{k} \sum_{\ell=1}^k \frac{\sum_{ij} \mathbf{W}_{ij} \delta(\mathbf{c}_i, \mathbf{c}_\ell) (1 - \delta(\mathbf{c}_j, \mathbf{c}_\ell))}{\sum_{ij} \mathbf{W}_{ij} \delta(\mathbf{c}_i, \mathbf{c}_\ell)}.$$

A cut between two clusters A and B is the sum of the weights between the clusters and the normalised cut is this sum divided by the sum of the weights of all edges incident to vertices in cluster A . Hence the k -way normalised cut is the mean normalised cut between each cluster and its complementary vertices, and the lower it is the better.

Figure 1 shows the k -way normalised cuts for both datasets. In both IASC and the Ning et al. method, since eigenvectors are recomputed every 10 iterations, this manifests itself as sudden changes in the k -way normalised cuts. These changes are more pronounced in the Ning et al. case. Except for the first few iterations, IASC matches or improves results on the exact approach, while keeping only 10% or fewer of the eigenvectors. The Ning et al. approximation leads to cuts worse than those obtained with the exact approach, except for the first few iterations immediately after exact recomputations of the eigenvectors. For the Bemol graphs Ning et al. failed to complete in a reasonable time frame, however on a smaller dataset of 1000 vertices, it resulted in cuts considerably worse than either IASC or the exact spectral approach.

6 Discussion

We have presented a novel incremental method for spectral graph clustering which updates the eigenvectors of the normalised Laplacian in a computationally efficient way. Such an algorithm is useful for finding clusterings in time evolving graphs such as biological and social networks, and the Internet. We compare our approach empirically to a naive spectral clustering approach which recomputes eigenvectors for each graph, and the method of Ning et al. and demonstrate improvements in the clustering on HIV and product information datasets as well as good cluster stability.

Recall that the approximation error of our proposed approach is kept under control by recomputing the exact eigen-decomposition after every R iterations. One promising extension to this work is to estimate this approximation error. This would allow one to limit the frequency of this computationally expensive step and offer a way to detect abrupt changes in the evolution of a graph sequence.

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